

# **Who is your neighbor?**

**--A brief introduction of Local Sensitive  
Hash and MinHash**

-- Xinzuo Wang

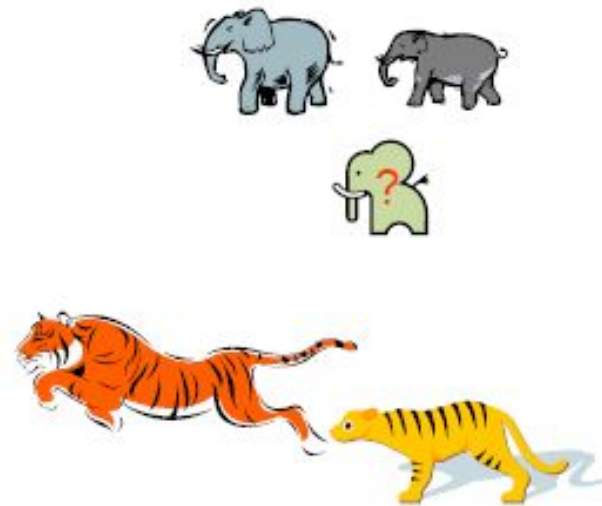
# An Outline:

- Near-optimal hashing algorithm for approximate near(est) neighbor problem
  - Local sensitive hash
- Finding similarity items from high-dimensional objects
  - MinHash

# Background

## Nearest Neighbor: Motivation

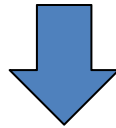
- Learning: nearest neighbor rule
- Database retrieval
- Vector quantization,  
also known as compression



# Background

Challenge : Curse of dimensionality

-- Difficult to compute large data set with large dimensions using original methods.



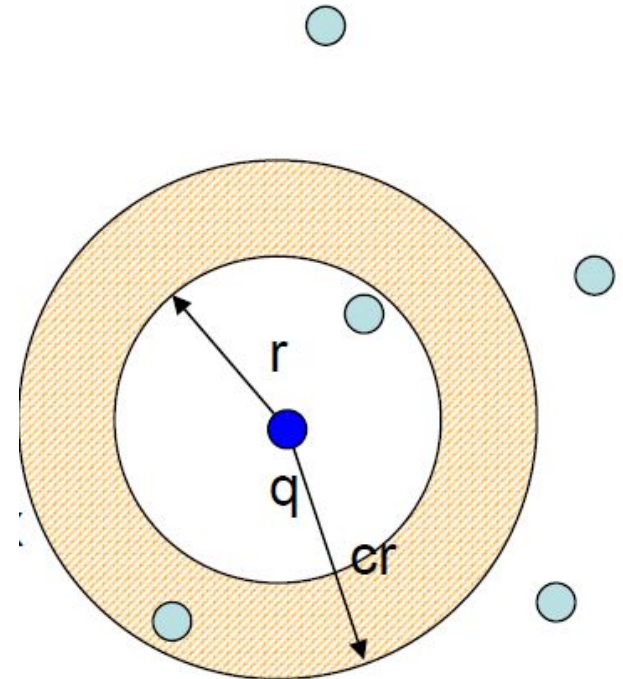
Approximate Near Neighbor

# Approximate Near Neighbor :

- $c$ -Approximate  $r$ -Near Neighbor: build data structure which, for any query  $q$ :
  - If there is a point  $p \in P$ ,  $\|p - q\| \leq r$
  - it returns  $p' \in P$ ,  $\|p' - q\| \leq cr$



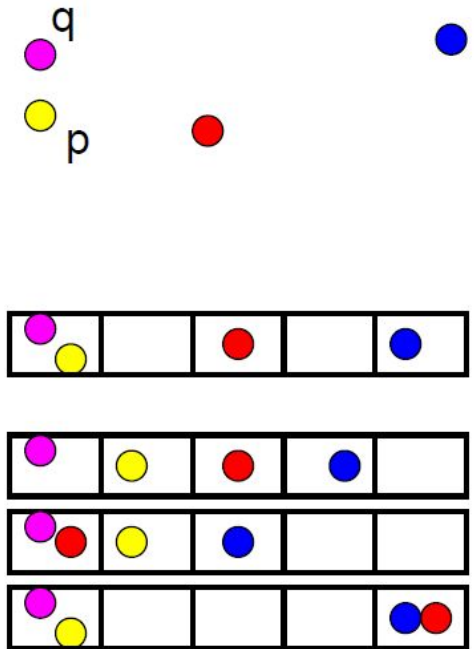
Here, *Local Sensitive Hash* is employed.



# Local Sensitive Hash :

-- *Transfer high dimensional data onto lower dimensions and preserve their 'distance'.*

- Idea: construct hash functions  $g: \mathbb{R}^d \rightarrow U$  such that for any points  $p, q$ :
  - If  $\|p-q\| \leq r$ , then  $P[g(p)=g(q)]$  is “high”
  - If  $\|p-q\| > cr$ , then  $P[g(p)=g(q)]$  is “small”
- LSH based on hamming norm.
- LSH based on projection.



# LSH on Hamming Norm:

Problem define :

Given a set  $P$ , return a point  $p \in P$ , for any query point  $q$  that  $d(p, q) < (1+\epsilon)d'(p, q)$ , where  $d'(p, q)$  is the smallest distance from  $q$  to  $p$

# LSH on Hamming Norm:

Hash functions:

$$h(p) = p_i, \text{ i.e., the } i\text{-th bit of } p$$

For example :

$$p = 10\mathbf{1}101010101$$

$$\text{if we define } i=3, h(p) = 1$$



# LSH on Hamming Norm:

The algorithm :

We use functions of the form  $g(p) = \langle h_1(p), h_2(p), \dots, h_k(p) \rangle$

- Preprocessing:
  - Select  $g_1 \dots g_L$
  - For all  $p \in P$ , hash  $p$  to buckets  $g_1(p) \dots g_L(p)$
- Query:
  - Retrieve the points from buckets  $g_1(q), g_2(q), \dots$ , until
    - Either the points from all  $L$  buckets have been retrieved, or
    - Total number of points retrieved exceeds  $2L$
  - Answer the query based on the retrieved points

# LSH on Hamming Norm:

- Some detailed analysis :
  - A crucial point : How to choose  $k$  and  $L$  ?

Intuitively, when  $k$  decreases, the precise of hash will increase;  
when  $L$  decreases, some points may be ‘unfortunately’ missed.

# LSH on Hamming Norm:

- If  $\|p-q\| \leq r_1$ , then  $P\{g(p)=g(q)\} \geq p_1$
- If  $\|p-q\| > r_2$ , then  $P\{g(p)=g(q)\} \leq p_2$  where  $(p_1 > p_2)$

Then we call the hash function is  $(r_1, r_2, p_1, p_2)$  sensitive

➔  $h(p)$  is  $(r, (1+\epsilon)r, 1 - \frac{r}{d}, \frac{r(1+\epsilon)}{d})$  sensitive

Thus, by correctly choose  $k$  and  $L$ , we can hash the point in  $(1+\epsilon)d'(p, q)$  to the same bucket with  $q$  with the probability more than  $p_1$

It can be proved that the optimal value of  $k$  is  $\log_{\frac{1}{p_2}} \left(\frac{n}{B}\right)$  and  $L$  is

$\left(\frac{n}{B}\right)^\rho$  where  $\rho = \frac{\ln(1/p_1)}{\ln(1/p_2)}$  and  $B$  is the max number of buckets.

# LSH on p-Norm – Based on projection

- Stable distribution :

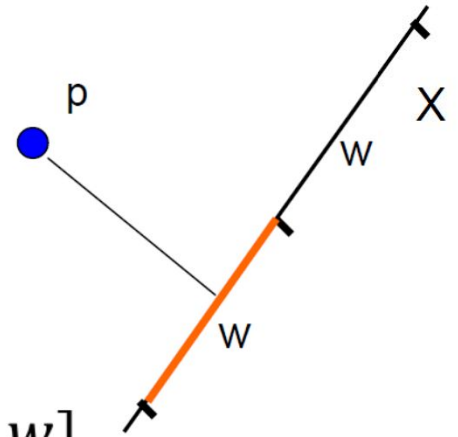
A stable distribution is called s-stable, if there exists  $p \geq 0$  such that for any real numbers  $v_1, v_2, v_3 \dots v_n$  and i.d.d. variables  $X_1, X_2, X_3 \dots X_n$  with distribution  $D$ , the random variable  $(\sum_i |v_i|^p)^{1/p} X$ , where  $X$  is a random variable with distribution  $D$ .

➡ One of the 2-stable distribution is the Normal distribution.

# LSH on p-Norm – Based on projection

- Hash function :

$$h_{\vec{a},b}(\vec{v}) = \left\lfloor \frac{\vec{a} * \vec{v} + b}{w} \right\rfloor \quad (R^d \rightarrow N)$$



$b$  is a real number chosen randomly from  $[0, w]$ ,  
 $\vec{a}$  is a  $d$ -dimension vector with entries chosen independently from a  $s$ -stable distribution.

# MinHash :

-- *Find similar items from high-dimensional objects.*

Main idea:

- Use *Min-Hash Signature* to preserve the ‘distance’ or ‘similarity’ between objects.
- Use *LSH* to find approximate nearest neighbors of each objects from *Min-Hash Signature*

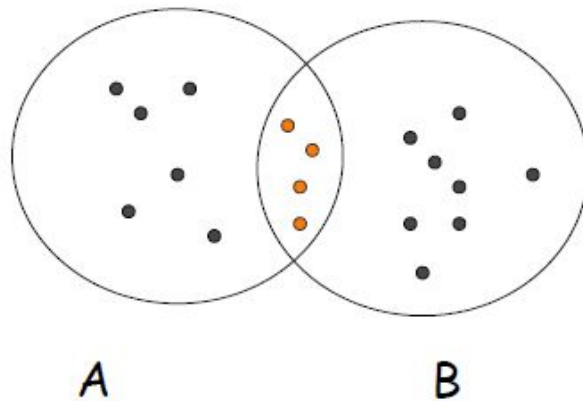
# MinHash :

- Define similarity :

Similarity metric, “distance”, for sets

Jaccard similarity:

$$\text{SIM}(A, B) := \frac{A \cap B}{A \cup B}$$



4 common  
18 total

$$\begin{aligned} \text{SIM}(A, B) &= 4/18 \\ &= 2/9 \end{aligned}$$

# MinHash :

- Hash functions :
- Minhash( $\pi$ ) of a set is the number of the row (element) with first non-zero in the permuted order  $\pi$ .

Element num	Set1	Set2	Set3	Set4
1	0	0	1	0
4	0	0	1	0
0	1	0	0	1
3	1	0	1	1
2	0	1	0	1
...				

$\Pi=(1,4,0,3,2)$



# MinHash -- Min-Hash Signature

$$\text{SIM}(Q, T) = P \{h(Q) = h(T)\}$$

How to compute  $P \{h(Q) = h(T)\}$  ?

- MinHash Signature :
  - Let  $h_1, h_2, \dots, h_n$  be different MinHash functions.
  - Then signature for set  $S$  is:  $\text{SIG}(S) = [h_1(S), h_2(S), \dots, h_n(S)]$

$\text{SIM}(Q, T) \approx$  the ratio of the ratio of equal elements of  $\text{SIG}(Q)$  and  $\text{SIG}(T)$

# An Outline:

- For example : A matrix of four sets with n elements

Element num	Set1	Set2	Set3	Set4	$x + 1$ mod 5	$3x + 1$ mod 5
0	1	0	0	1	1	1
1	0	0	1	0	2	4
2	0	1	0	1	3	2
3	1	0	1	1	4	0
4	0	0	1	0	0	3
...						

MinHash signature :

	set1	set2	set3	set4
1	1	3	2	1
2	1	3	2	1
3	1	2	4	1

# LSH based on MinHash

Mean idea:

- Divide the signature matrix rows into  $\mathbf{b}$  bands of  $\mathbf{r}$  rows
- Hash the columns in each band with a basic hash-function, each band divided to buckets.

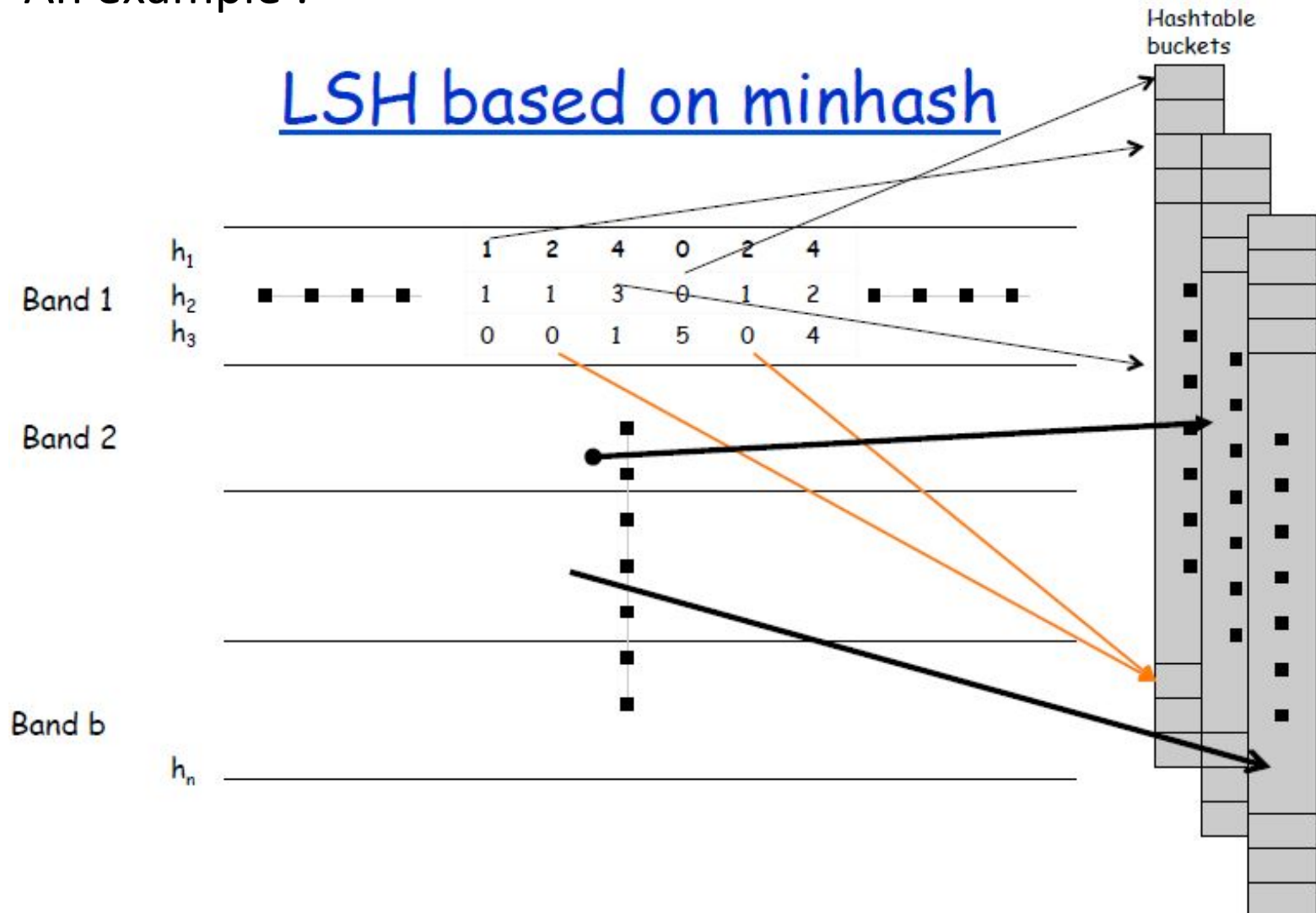
If sets  $S$  and  $T$  have same values in a band, they will be hashed

into the same bucket in that band.

- For nearest-neighbor, the candidates are the items in the same bucket as query item, in each band.

# LSH based on MinHash

- An example :



# LSH based on MinHash

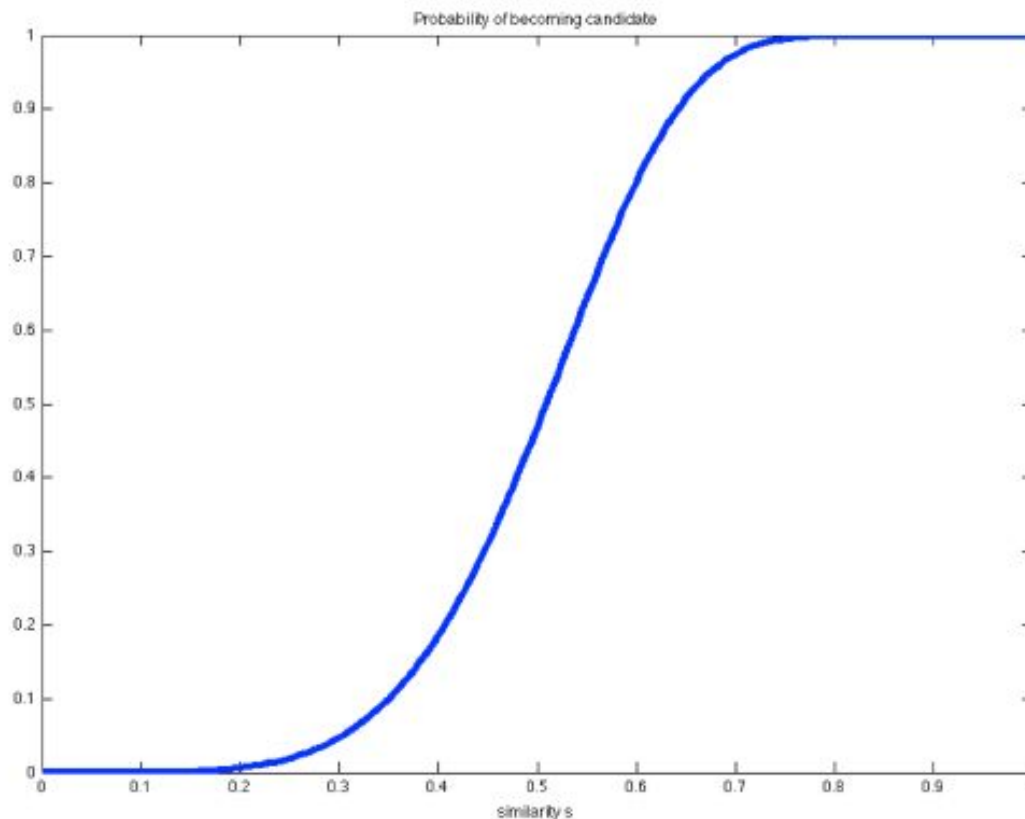
- Analysis :
  - What is the probability the S will be chosen in the same bucket with T when given  $\text{SIM}(S, T) = s$ ?
- $P\{Q \text{ and } T \text{ agree on all rows in a band}\} = s^r$
- $P\{S \text{ will be chosen as the candidate}\}$ 
  - $= P\{S \text{ and } T \text{ agree at least on one band}\}$
  - $= 1 - (1 - s^r)^b$

# LSH based on MinHash

- Relationship between  $\text{SIM}(S, T)$  and the probability of becoming candidate with the value of  $r$  and  $b$ .

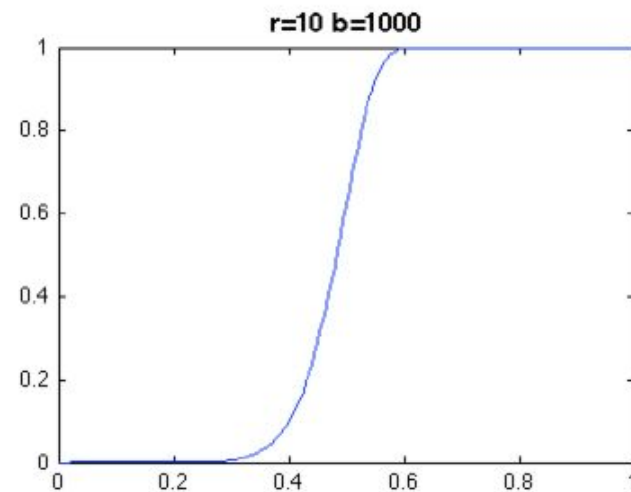
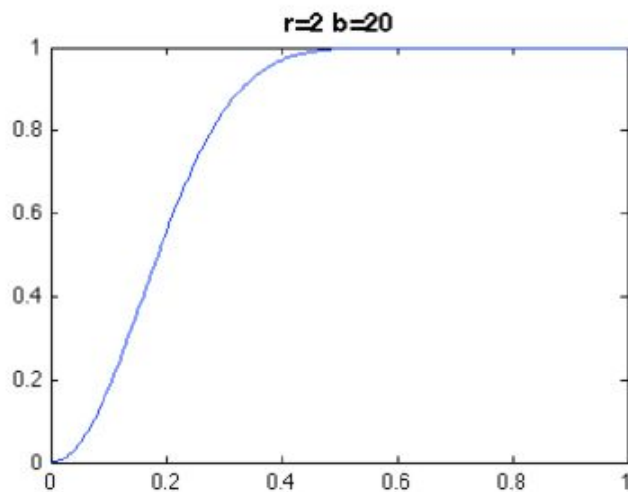
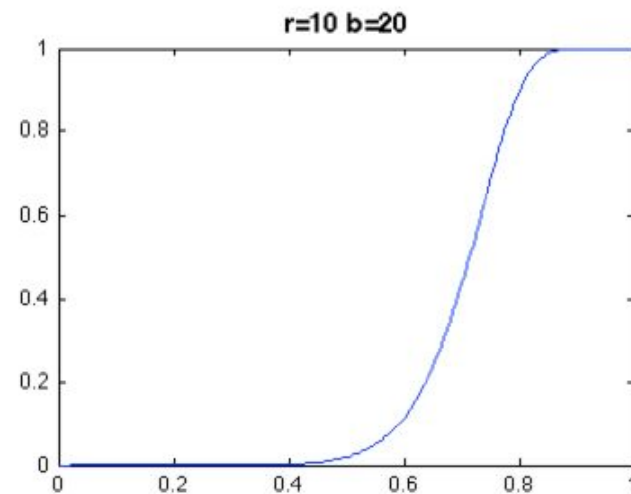
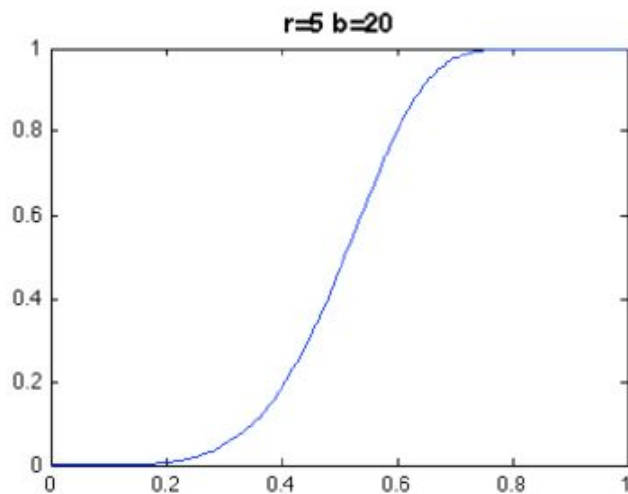
S-curve

$r = 5$   
 $b = 20$



# LSH based on MinHash

- When choosing different  $r$  and  $b$  :



# Summary :

- Local sensitive hash -- Find nearest neighbors
  - Transform high-dimensional data onto lower dimensions with similarity preserved(similar points are more likely to be hashed in the same bucket).
- MinHash – Find similarity items
  - Use MinHash Signature to approximately preserve similarities between each other (aims to reduce data size).
  - Use LHS to find neighbors of each node.



# References :

- Website [1] lists a series of algorithm of LSH and the implement of Euclid Norm.
- [2] is a brief introduction of LSH.
- Chapter 3.3 in [3] introduced MinHash algorithm.
- Monograph [4] systematically introduced Nearest-neighbor methods.

[1] <http://www.mit.edu/~andoni/LSH/>

[2] <http://web.iitd.ac.in/~sumeet/Slaney2008-LSHTutorial.pdf>

[3] Rajaraman A, Ullman J D. Mining of massive datasets[M]. Cambridge University Press, 2011.

[4] Shakhnarovich G, Indyk P, Darrell T. Nearest-neighbor methods in learning and vision: theory and practice[M]. 2006.

Thanks !!